

CS32 Discussion Week 3

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Outline

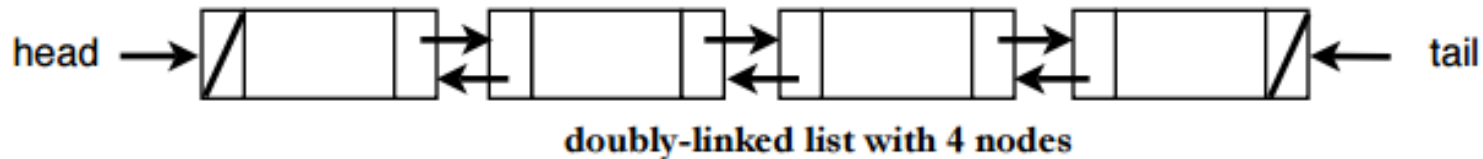
- Doubly Linked List
- Sorted Linked List
- Reverse a Linked List

Doubly Linked List

- A linked list where each node has two pointers:
 - Next – pointing to the next node
 - Prev – pointing to the previous node
- struct Node {
 - int value;
 - Node *Next;
 - Node *Prev;
- };

Doubly Linked List

- That's how it looks like:



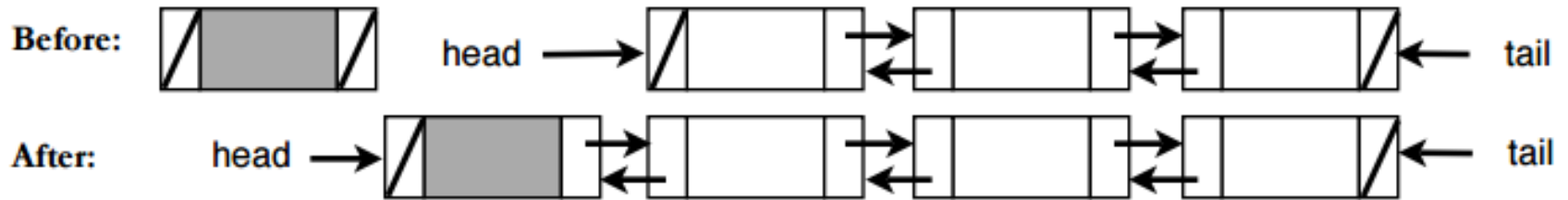
- Features to capture a DLL:
 - Two pointers: *head*, *tail*
 - $\text{head} \rightarrow \text{prev} = \text{NULL}$
 - $\text{tail} \rightarrow \text{next} = \text{NULL}$
 - $\text{head} == \text{tail} == \text{NULL}$ when list is empty

Insertion

Four different conditions to insert a new node P

1. Insert before the head;
2. Insert after the tail;
3. Insert somewhere in the middle
4. When list is empty;

Insertion (Before head)



- 1) Set the *prev* of *head* to the new node *p*
 - `head -> prev = p;`
- 2) Set the *next* of *p* to *head*
 - `p -> next = head;`
- 3) *p* becomes the new *head*
 - `head = p;`
- 4) `head -> prev = NULL;`

Insertion (after tail)

- Quite the same as insertion before head:

```
tail -> next = p;  
p -> prev = tail;  
tail = p;  
p -> next = NULL;
```

Insertion in the middle (*after node q*)

- 1) Fix the next node of q first:
 - Node $*r = q \rightarrow \text{next}$;
- 2) Point both *next of q* and *prev of r* to p
 - $q \rightarrow \text{next} = r \rightarrow \text{prev} = p$;
- 3) Point both sides of p to q and r respectively:
 - $p \rightarrow \text{prev} = q$;
 - $p \rightarrow \text{next} = r$;

Insertion in the middle (*after node q*)

- Or do it without r:
 - $p \rightarrow \text{prev} = q;$
 - $p \rightarrow \text{next} = q \rightarrow \text{next};$
 - $q \rightarrow \text{next} = q \rightarrow \text{next} \rightarrow \text{prev} = p;$

Insertion (to an empty list)

- How do we represent an empty list?
 - `head == NULL` (Or `tail == NULL`; Or `head == tail == NULL`)
- 1) Insertion, just set p as *head* (as well as *tail*):
 - `head = tail = p;`
- 2) Don't forget to set `NULL` on both sides:
 - `p->next = p->prev = NULL;`

Search

- Just like the singly linked list.

```
Node* Search(int key, Node* head){
    Node *q = head;
    while(q != NULL) {
        if(q -> value == key) return q;
        else q = q -> next; //iterate to
the next node
    }
    return NULL;
}
```

```
Node* Search(int key, Node* tail){
    Node *q = tail;
    while(q != NULL) {
        if(q -> value == key) return q;
        else q = q -> prev; //iterate to
the previous node
    }
    return NULL;
}
```

Removal

- More complex than singly linked list.
 - Check if the node p is the head ($p == \text{head}$). Let this boolean be **A**.
 - Check if the node is the tail ($p == \text{tail}$). Let this boolean be **B**.

Removal

- Four cases:
 - **Case 1 (A, but not B):** P is the head of the list, and there is more than one node.
 - **Case 2 (B, but not A):** P is the tail of the list, and there is more than one node.
 - **Case 3 (A and B):** P is the only node.
 - **Case 4 (not A and not B):** P is in the middle of the list.

Removal Case 1 (P is head)

- 1) Update *head*
 - head = head -> next;
- 2) delete p
 - delete p ;
- 3) Set the *prev* of *head* to NULL
 - head -> prev = NULL;

Removal Case 2 (P is tail)

- 1) Update *tail*
 - $\text{tail} = \text{tail} \rightarrow \text{prev};$
- 2) delete p
- 3) Set the *next* of *tail* to NULL
 - $\text{tail} \rightarrow \text{next} = \text{NULL};$

Removal Case 3 (P is the only node)

- 1) Empty the linked list:
 - head = tail = NULL;
- 2) delete p :

Removal Case 4 (P is in the middle)

- 1) Fix the *prev* and *next* of p :
 - Node $*q = p \rightarrow \text{prev}$;
 - Node $*r = p \rightarrow \text{next}$;
- 2) Concatenate q and r :
 - $q \rightarrow \text{next} = r$;
 - $r \rightarrow \text{prev} = q$;
- 3) Delete p

Removal Case 4 (Equivalent implementation)

- If we do not fix with q and r :
 - $p \rightarrow \text{prev} \rightarrow \text{next} = p \rightarrow \text{next};$
 - $p \rightarrow \text{next} \rightarrow \text{prev} = p \rightarrow \text{prev};$
 - delete p ;

Removal summary

```
void removeNodeInDLL(Node *p, Node& *head, Node& *tail) {
    if (p == head && p == tail) //case 3
        head = tail = NULL;
    else if (p == head) { //case 1
        head = head -> next;
        head -> prev = NULL;
    }
    else if (p == tail) { //case 2
        tail = tail -> prev;
        tail -> next = NULL;
    }
    else { //case 4
        p -> prev -> next = p -> next;
        p -> next -> prev = p -> prev;
    }
    delete p;
}
```

Copying a doubly linked list

- 1) Create *head* and *tail* for the new list
- 2) Iterate through the old list. For each node, **copy its value to a new node.**
- 3) Insert the new node to the tail of the new list.
- 4) Repeat 3 until we have iterated the entire old list. Set NULL before *head* and next of *tail*.

Copy a Doubly Linked List

```
void copyDDL(Node *head_o, Node *tail_o, Node& *head_n, Node& *tail_n) {
    if (tail_o == NULL) { //the original list is empty
        head_n = tail_n = NULL; return;
    }
    Node *q = head_o; //iterator
    Node *p = new Node();
    p -> value = q -> value;
    head_n = tail_n = p;
    q = q -> next;
    while (q) {
        p = new Node();
        p -> value = q -> value;
        tail_n -> next = p;
        p -> prev = tail_n;
        tail_n = tail_n -> next;
        q = q -> next;
    }
    head_n -> prev = tail_n -> next = NULL
}
```

} insertion for the first node is different

} Copy value to the new node

} Append the new node to the tail of the new list, and update tail.

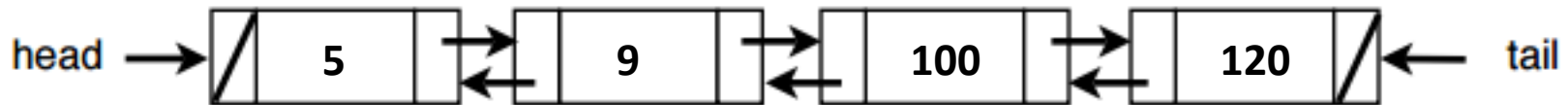
Cautions about coding with a linked list

- To draw diagrams of nodes will be extremely helpful.
- When copying a linked list, only copy stored values to new nodes. Do not copy pointers.

Sorted linked list

Do we need to search the entire linked?

- Consider an ascending sorted (doubly linked) list:



- Do we need to search through all the nodes when we're searching for:
 - 8
 - 50
- A way to optimize the search: sorted list and early stop

How to implement an ascending sorted (doubly linked) list?

- Change the insertion. Find the node q whose value is the *greatest lower bound* to the new node p .
- 1) Check if $head$'s value is larger than p 's. If so, insert p as head.

```
if (p -> value < head -> value) {  
    p -> next = head;  
    head -> prev = p;  
    p -> prev = NULL;  
    head = p;  
}
```

How to implement an ascending sorted (doubly linked) list?

- 2) if not, iterate through the list until we find the node q whose value is the *greatest lower bound* to p .
 - Note: $q \rightarrow \text{next}$ can either have a value larger than q , or is NULL.

```
Node *q = head -> next;
```

```
while (q && q -> value < p -> value) {
```

```
    if (q -> next == NULL || q -> next -> value > p -> value) break;
```

```
    q = q -> next;
```

```
}
```

How to implement an ascending sorted (doubly linked) list?

- 3) Insert p after q :

```
if (q -> next != NULL)
```

```
    q -> next -> prev = p;
```

```
p -> next = q -> next; /* if q is the last node then its next is  
already NULL */
```

```
p -> prev = q;
```

```
q -> next = p;
```

Insert into an ascending doubly linked list

```
void insert(Node *p, Node *head) {
    if (p -> value < head -> value) {
        p -> next = head;
        head -> prev = p;
        p -> prev = NULL;
        head = p;
        return;
    }
    Node *q = head -> next;
    while (q -> value < p -> value) {
        if (q -> next == NULL || q -> next -> value > p -> value) break;
        q = q -> next;
    }
    if (q -> next != NULL)
        q -> next -> prev = p;
    p -> next = q -> next;
    p -> prev = q;
    q -> next = p;
}
```

Early stop in searching a sorted linked list

We stop the iteration once we see a node which stores a value that is larger than key.

```
Node* Search(int key, Node* head){
    Node *q = head;
    while(q != NULL && q -> value <= key) {
        if(q -> value == key) return q;
        else q = q -> next; //iterate to the next node
    }
    return NULL;
}
```

What about removal and update?

Why *early stop* technique saves cost

- Spare cost from search to insertion and update
- However, search is called massively, but insertion and updates are not.
- $O(n/2)$ cost is saved for each search (assuming the data complies with a uniform distribution)

Reverse Linked List (Leetcode #206, easy)

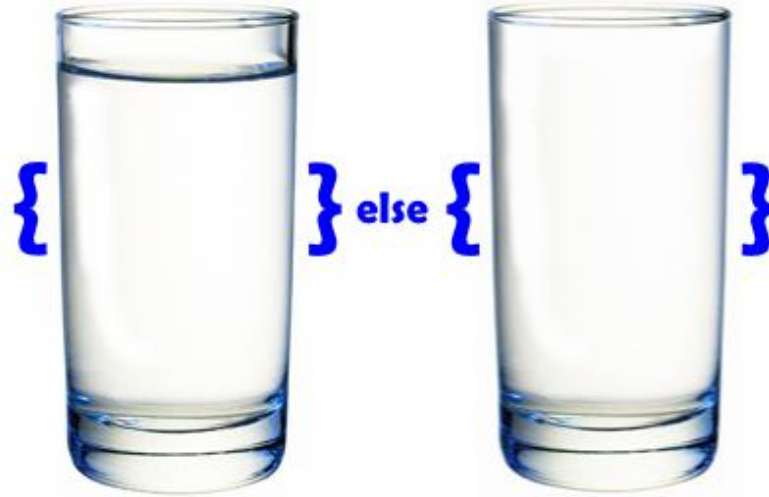
Given a **singly** linked list, reverse every node of it (i.e. each next points to the previous node).

```
Node* reverseList(Node* head)
```


Solution

```
Node* reverseList(struct ListNode* head) {  
    Node *prev=NULL,*cur=head,*next;  
    while(cur) {  
        next = cur->next;  
        cur->next = prev;  
        prev = cur;  
        cur = next; }  
    return prev;  
}
```

if (\$thirsty==TRUE)



Bugs in your software are actually special features :)